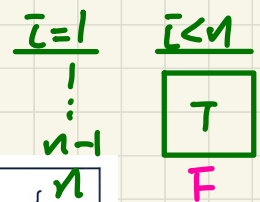


Friday, February 3

Written Test 1 Review

* KCN only executed before outer loop execs (when $j=n$)



Count # of Primitive Operations

```

1 int sumMaxAndCrossProducts (int[] a, int n) {
2   int max = a[0]; 2
3   for(int i = 1; i < n; i++) {
4     if (a[i] > max) { max = a[i]; }
5   }
6   int sum = max; 1
7   for (int j = 0; j < n; j++) {
8     for (int k = 0; k < n; k++) {
9       sum += a[j] * a[k];
10    }
11  }
12  return sum;

```

1. # times $i < n$ evaluated? n
 2. # times body of loop exec? $n-1$

$$1 + n + 2 \cdot (n-1) + 4 \cdot (n-1) = 7n - 5$$

for each value of j making $j < n$ true, $k < n$ could be evaluated $(n+1)$ times.

j	k	1	2	...	n-1	n
0	0	1	2	...	n-1	n
1	0	1	2	...	n-1	n
2	0	1	2	...	n-1	n
...
n-1	0	1	2	...	n-1	n

$$1 + n + (n+1) + n \cdot (n+1) + \dots$$

$$2 \cdot n + 2 \cdot n^2 + 5 \cdot n^2 + \dots = ?$$

$j < n \rightarrow (n)$

Count # of Pos

$$24 + 21 + 18 + 15 + 12$$

$$= \frac{(24 + 12) * 5}{2}$$

$$(n+1) + n + (n-1) + \dots + 2 + 1$$

$$= \frac{(n+1) + 2 * n}{2}$$

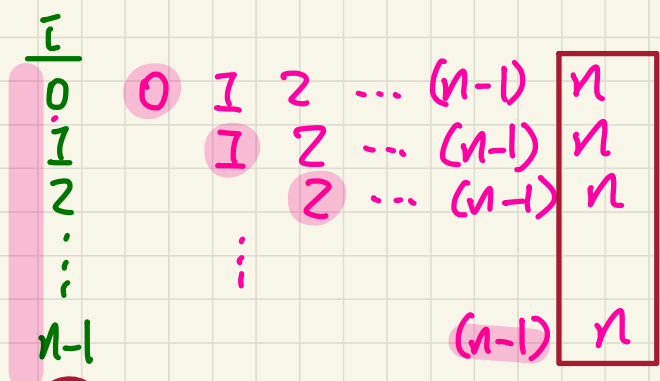
```

1 int triangularSum (int[] a, int n) {
2   int sum = 0;
3   for (int i = 0; i < n; i++) {
4     for (int j = i; j < n; j++) {
5       sum += a[j];
6     }
  }
  return sum;
}

```

how many terms?
 $n + (n-1) + \dots + 1$
 $= \frac{(n+1) * n}{2}$

when i is between 0 and $n-1$,
 $j < n$ is evaluated between i and n .



$j < n \rightarrow \text{E}$

$$1 + n + \frac{n+1}{i < n} + \frac{(n+1+2) * n}{2}$$

$j < n$

$$+ \frac{n * 2}{i++} + \frac{(n+1) * n}{2} * 2 + \frac{(n+1) * n}{2} * 3$$

$j++$ $C5$

$$+ 1 = ?$$

$i < n$
 \downarrow
 E

```
String[] insertAt(String[] a, int n, String e, int i)
String[] result = new String[n + 1];
for(int j = 0; j <= i - 1; j++){ result[j] = a[j]; }
result[i] = e;
for(int j = i + 1; j <= n; j++){ result[j] = a[j-1]; }
return result;
```

for (int j = 0; j <= i-1; j++) {
 for (int k = i+1; k <= n; k++) {

<u>j</u>	<u>k</u>	
0	<u>i+1</u> <u>i+2</u> ... n	$n - (i+1) + 1$ $= \boxed{n - i}$ $(n - i) \cdot i$ $= \underline{n \cdot i} - \textcircled{i^2}$ <p style="text-align: right; margin-right: 50px;">Constant</p>
1	<u>i+1</u> <u>i+2</u> ... n	
⋮	⋮	
<u>i-1</u>	<u>i+1</u> <u>i+2</u> ... n	

} }
}

O(n)

```

int count = 0;
for (int i = n/2; i <= n; i++)
    for (int j = 1; j + n/2 <= n; j = j++)
        for (int k = 1; k <= n; k = k * 2)
            count++;

```

$j + \frac{n}{2} \leq n$
 $j \leq n - \frac{n}{2} = \frac{n}{2}$

Assume $n = 1000$

- $k = 1 = 2^0$
- $2 = 2^1$
- $4 = 2^2$
- $8 = 2^3$
- \vdots
- $512 = 2^9$

$10 = \lceil \log_2 1000 \rceil$

How many times j changes its value?

$\approx \frac{n}{2}$

i	j	k	\dots	$\frac{n}{2}$
$\frac{n}{2}$	1	1	$\dots \log n$	$\frac{n}{2}$
$\frac{n}{2} + 1$	1	2	3	$\frac{n}{2}$
$\frac{n}{2} + 2$	1	2	3	$\frac{n}{2}$
\vdots				
n	1	2	3	$\frac{n}{2}$

$O\left(\frac{n}{2} \cdot \frac{n}{2} \cdot \log n\right) = O\left(\frac{n^2}{2} \cdot \log n\right)$